

# Solving the Traveling Salesman Problem using Genetic Algorithms

Luis F. Copertari

**Abstract—** The Traveling Salesman Problem (TSP) is solved using genetic algorithms. Two alternative ways to represent the solution (genes in series and genes in two columns) are used and different mutation rates as well as different population sizes are considered. Two computer programs are created to solve the problem and yield experimental data. The data is analyzed and discussed and conclusions are derived. An equation relating the percentage of the population with the same solution ( $w$ ) and the mutation rate ( $g$ ) is discussed. This equation and the tradeoff between population size, mutation rate and chromosome portraying method are the theoretical contributions of this work, as well as an innovative way to solve the TSP, which is a hard problem to solve using traditional optimization, such as linear programming.

**Index Terms—** Traveling Salesman Problem, Genetic Algorithms, Evolutionary Computation

## I. INTRODUCTION

The Traveling Salesman Problem (TSP) is a very well-known problem in operations research, in which one salesman tries to visit  $n$  cities in a given tour. The problem is to find the sequence of visits that minimizes the sum of all of the distances between cities [1]. Checking all possible tours requires doing  $n!$  additions. For a sample problem with 27 cities, that means  $1.088886945 \times 10^{28}$  sums. A modern computer is capable of approximately one billion operations per second [2]. That means it would take 345,283,785,200 years to cover all possibilities a brute force approach. An impossible problem!

This problem can be solved using linear programming [3], but for large values of  $n$  it takes a relatively long time (hours, days, months, or even years) to come up with a solution. The approach is to use genetic algorithms. Although genetic algorithms do not guarantee an optimal solution, they provide good solutions and sometimes even the optimal solution in little time. Moreover, a solution (although maybe not the optimal) is always available. It is not necessary to wait until the program finishes. The user may request the best solution found so far at any time.

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Genetic algorithms rely on evolution to find a solution. Each individual in the population represents one possible solution. The individuals are represented so that it is possible to create new generations by reproducing the best fitted individuals of the previous generation.

Two alternative ways to represent TSP tours were considered. The first one is by arranging the cities to visit in a row. The second one is to arrange the cities in two columns. Each approach entails a different way of reproducing the population. Both approaches were tried with different mutation rates to see which approach leads faster to the solution or provides the optimal solution.

First, the methodology is discussed. Second, the theory on the TSP and genetic algorithms is reviewed. Third, the experimental results are considered. Finally, the findings are discussed and conclusions are derived.

## II. METHODOLOGY

The methodology is basically the scientific method. The scientific method is a series of steps, which are iterated in a series of successive and ever increasingly closer approximations to the truth. A theoretical approach had to come up based on the observation of the problem under study, test the theory in the laboratory in a series of experiments, and if the results were not within theoretical parameters, there was the need to go back to the theory and try to be more accurate based on the available observations.

Throughout the years, the scientific method has been discussed thoroughly. The specific approach used is based on the ideas presented by Gauch Jr. [4] and Wilson [5], which consists of the following steps:

1. Observation. The direct observation from reality is the basis of all discoveries. The problem was first considered and then it was realized that there are at least two different ways to portray the information: genes in series and genes in two columns.
2. Hypothesis. The right combination of mutation rate, population size, and model (genes in series or genes in columns) are important to determine how good the solution is and how fast the solution is reached.
3. Theory. What is to know about the TSP was considered as well as its implementation using genetic algorithms.
4. Experiment. Two computer programs using Delphi were created to solve the two genetic algorithm problems and obtain data for each representation, each possible mutation rate and each population size.
5. Conclusions. After analyzing the data obtained in the laboratory, the results were considered and discussed. Finally, conclusions are derived.

### III. THEORY

#### A. The Traveling Salesman Problem

The Traveling Salesman Problem (TSP) is a classical problem in operations research. There are  $n$  cities, and it is supposed a salesman (or as an extension a load, a truck, or a plane) has to visit each and every city in order to minimize the total distance travelled.

The TSP is a NP problem, that is, the number of steps required to solve it grows exponentially as the problem size grows. The above means that to solve the TSP for typical examples from real life, where there is a fleet of hundreds of Wonder or Coke delivery trucks, or any other case that can be thought of, requires a prohibitive number of steps to obtain the solution by solving using linear programming, for example. That is, the computer would take too much time (years maybe) to solve a problem large enough so that it matters in real life.

Table 1 shows the distances between  $n$  cities. Notice that it is required to specify the distance between city 1 and city  $j$ , where  $j$  goes from 2 to  $n$ , which is done in the first row. For  $i=1$  and  $j=1$ , there is a distance of zero, because the distance between city one and city one is none. However, to avoid going from one city to the same one, a very large distance is assigned, such as 100,000 (for all  $i = j$ ). In the second row, there is the distance between city 2 and city  $j$  (where  $j$  goes from 3 to  $n$ ; notice that  $j=1$  is not included, because the distance between city 2 and city 1 is already in the first row as the distance between city 1 and city 2). And so on. In general, the distance between city  $i$  and city  $j$ , where  $j$  goes from  $i+1$  to  $n$ , is denoted as  $d_{ij}$ . Notice that in this way the upper triangle of distances is obtained, marked in grey, which is duplicated in the lower triangle of distances. However, it is possible in reality that the distance to go from city  $i$  to city  $j$  ( $d_{ij}$ ) may be different than the distance to go from city  $j$  to city  $i$  ( $d_{ji}$ ).

**Table 1.** Distances between cities for the TSP.

City\City	1	2	3	...	n
1	100,00 0	$d_{12}$	$d_{13}$	...	$d_{1n}$
2	$d_{21}=d_{12}$	100,00 0	$d_{23}$	...	$d_{2n}$
3	$d_{31}=d_{13}$	$d_{32}=d_{23}$	100,00 0	...	$d_{3n}$
...	...	...	...	...	...
n	$d_{n1}=d_{1n}$	$d_{n2}=d_{2n}$	$d_{n3}=d_{3n}$	...	100,00 0

#### B. Genetic Algorithms

Genetic algorithms were developed by John Holland [6] during the 1960s and 1970s and popularized by one of his students, David Goldberg [7]. Genetic algorithms rely on the principles of genetics and natural selection to find solutions. De Jong [8] showed the usefulness of genetic algorithms for function optimization and tried to find optimal parameters for genetic algorithms.

Some of the advantages of genetic algorithms, taken from Haupt & Haupt [9], are:

- Optimizes with continuous or discrete variables.

- Doesn't require derivative information.
- Simultaneously searches from a wide sampling of the cost surface.
- Deals with a large number of variables.
- Is well suited for parallel computers.
- Optimizes variables with extremely complex cost surfaces (they can jump out of a local minimum).
- Provides a list of optimum variables, not just a single solution.
- May encode the variables so that the optimization is done with the encoded variables.
- Works with numerically generated data, experimental data, or analytical functions.

Genetic algorithms are very intriguing and they produce stunning results when traditional optimization techniques, such as linear programming, fail miserably. Sometimes, such as when evaluating an analytical function, calculus may outperform genetic algorithms. Also, the fact that a large population of solutions is required means that traditional methods may find the solution faster. However, if a parallel computer is available, genetic algorithms can take advantage of that because each processor may calculate the fitness of a group of individuals. Genetic algorithms are ideally suited for parallel computing.

There are five steps in the genetic algorithm process:

1. Encoding (genesis): Each individual in the population is generated such that it portrays one possible solution in a given string and it is possible to do the crossover using such encoding scheme.
2. Evaluation: Each solution of the individuals in the population is evaluated to find how well it fits the solution to the problem and each individual is assigned a score.
3. Crossover: Based on the score, only the best individuals are crossed or reproduced to create new individuals.
4. Mutation: In some cases, random changes or twitches are applied to some individuals randomly chosen for such purpose.
5. Decoding: After the process has finished, given so many generations have passed, the solution encoded in the representation of the individual with the best score is decoded and transformed into a solution to the problem. If an optimum (or at least a local optimum) has been found, a relatively large percentage of the population will have found such best solution.

#### C. Solving the TSP with Genetic Algorithms

It is possible to solve the traveling salesman problem using genetic algorithms, although Kirkpatrick, Gelatt and Vecchi [10] use simulated annealing to solve the problem. There are several ways to do the encoding. An encoding known as matrix representation, discussed by Homaifar, Guan and Liepins [11], as well as Michalewicz [12], is possible. In matrix representation, the tours are represented in a binary (0 and 1) matrix so that if there is a one in row  $i$ , column  $j$ , that means the salesman is going from city  $i$  to city  $j$ . This kind of matrix is also called a precedence matrix in operations research [1][3].

The matrix shown in Figure 1 indicates that the salesman is

going from city 1 (row 1) to city 2 (because in row 1, the one is in column 2). Also, that from city 2 goes to city 3 and from city 3 to city 1. As can be seen, the tour matrix from Figure 1 is valid because it constitutes a valid tour, that is, there are no repeated sequences such as, for example, having the possibility of going from city 1 to both cities 2 and 3.

Figure 1. Matrix representation using a precedence matrix.

$$\begin{matrix} & \begin{matrix} 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \left[ \begin{array}{ccc} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{array} \right] \end{matrix}$$

The problem with the matrix representation is the crossover operation. In crossover (or reproduction) the characteristics of the individual represented by their chromosome (encoding) are mixed to create new individuals. However, matrix representation as it is does not guarantee valid offspring.

The TSP problem can also be represented by a string of integers in two different ways. The first one, given by Homaifar et al. [11], Michalewicz [12], Goldberg [7] and Wroblewski [13], is using a string, such as  $a_1a_2a_3...a_n$ . The previous string means that the salesman would go from city  $a_1$  to city  $a_2$ , from  $a_2$  to  $a_3$ , and so on until city  $a_n$ , and also from city  $a_n$  to city  $a_1$  in order to complete the circuit.

For convenience, the string notation is going to be used. The encoding and decoding operations are very straightforward to do. However, the evaluation, crossover and mutation operations have to be explained. The evaluation operation is simply taking the distance between the first and second cities, adding the distance between the second and third city, and so on until city  $n$ , where the distance between city  $n$  and city 1 is added to complete the circuit.

The crossover operation is a little bit more complicated to explain. Let say there are strings with  $n$  cities. The best first half of individuals is taken and two of them are chosen at random from this pool of candidates. The first individual chosen is the male, the second the female. A random number between 0 and  $n$  is obtained. Let say such number is  $d$ . This random number indicates the breaking point in the male. If the number obtained is  $n$ , it means that all of the numbers in the male string must pass to the offspring (cloning of the male). If the number is zero, it means all the numbers in the female pass intact to the offspring (cloning of the female). If a number between 1 and  $n-1$  is obtained, both individuals have to be combined. The first  $d$  numbers from the male are taken, and then all the non-repeating numbers from the female are taken to create a new sequence of numbers in which the left side of the male is preserved.

Consider the example shown in Figure 2, where  $n=5$ . The top numbers in the male chromosome are the cutting positions. Zero is at the beginning and five at the end. Let say  $d = 2$ . The cutting point in the male is indicated at the top of Figure 2. All the numbers to the left of the cutting point of the

male are preserved in the offspring, so that the offspring chromosome begins with 5-3. To complete the sequence in the offspring, all the valid numbers from the female chromosome are taken. The first number in the female is 3, but 3 cannot be used because there is already a 3 from the male between position 1 and 2. The next number in the female is 4, and since 4 is not repeated, it is added. The third number in the female is 5, but that is not valid because the offspring already has a five. The fourth number in the female is 2, and since it is valid (not repeated) it is added. Finally, the fifth number in the female chromosome is 1, which is also valid and added to the offspring. In this way valid chromosomes as offspring can be generated.

Figure 2. Crossover operation for genes in series.

Male

$$\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 5 & - & 3 & | & 1 & - & 2 & - & 4 \end{matrix}$$

Female

$$\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 3 & - & 4 & - & 5 & - & 2 & - & 1 \end{matrix}$$

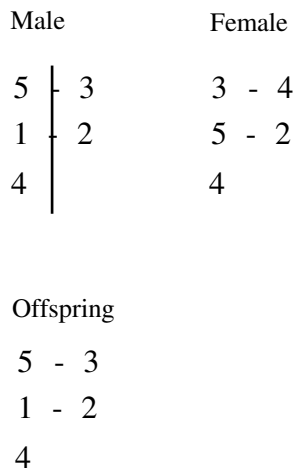
Offspring

$$\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 5 & - & 3 & - & 4 & - & 2 & - & 1 \end{matrix}$$

It is also important to notice that apparently similar chromosomes may refer to the same tour. Consider for example  $n=5$  and a tour such as 5-3-1-2-4. The chromosome 3-1-2-4-5 represents the same information as 5-3-1-2-4 because city 5 is at the end in the chromosome 3-1-2-4-5, but since the chromosomes are circular, after city 5 follows city 3, which is the same as in chromosome 5-3-1-2-4. The rest of the sequence after this is the same. Another variant is to invert the tours. Chromosome 4-2-1-3-5 is the same as chromosome 5-3-1-2-4 and chromosome 3-1-2-4-5, because the tour is still the same even though the order has been reversed.

Another variant being considered is to arrange cities on columns instead of series. This new arrangement changes the way in which the crossover occurs. Figure 3 shows the same case as the one in Figure 2. Notice that now, genes are arranged in two columns. For reproduction, the left side of the male is taken (5-1-4) and the right side is completed using the sequence from the female chromosome. In this case, the first city is 3, and since there is no 3 in 5-1-4, it is added. The second city is 4, but it is not included since 5-1-4 already contains 4. The third gene is 5, which is also not included since 5 is already in the sequence of the left column in the offspring (5-1-4). Then comes 2, which is included for not being in the left side sequence. Finally, there is 1, which is not included (the chromosome has been completed anyway), because it is already in the left side of the male chromosome (5-1-4).

Figure 3. Crossover operation for genes in columns.



Mutation is very simple to perform. It is just a matter of drawing a random number for each offspring and if it is less than or equal to a given mutation percentage rate,  $w$ , then such individual is to mutate. To perform the mutation on any given individual, two random numbers are simply drawn between 1 and  $n$  and swap the cities in those positions. Notice that now the positions do not refer to the space between the cities but to the space of the cities themselves.

There is a limit to mutation. Let  $g$  be the percentage of the population required to have the same solution in order to finish the run. Also, let  $w$  be the mutation rate, that is, the percentage of the population that is expected to mutate in a given iteration. Assume that for a particular generation there is, before mutation, 100% agreement in the solution and that  $g = 0.70$ . Since  $100\% > 70\%$ , if there were no mutation, the run would be finished. Also, let say that the mutation rate is 50%. Since out of the 100%, 50% of the population is expected to mutate, we would end up having only 50% of the population untouched with the original solution. In this situation, the algorithm would never converge. In general, equation (1) applies.

$$g+w < 1 \tag{1}$$

Equation (1) indicates that in order for the algorithm to have a chance to converge,  $g+w$  must be less than one. A value for  $g$  of 0.7 (70%) is used. Clearly, a mutation rate of 50% ( $w=0.5$ ) is too high and the algorithm would never converge under such conditions since  $0.7+0.5 = 1.2 > 1$ .

#### IV. TESTING

To test whether or not a given encoding scheme (it could be genes in series or genes in columns) implies reaching the solution faster or assuring an optimal solution, a number of  $r$  tests are going to be conducted. One of the variables in the tests is the encoding mechanism: series or columns. Another

variable is the mutation rate: 5%, 10%, 20% or 50%.

There is also a final variable. Considering populations throughout human history, the question concerning the size of the population having an impact in the speed in which the population evolved arises. In the distant past, when our ancestors were in the stone-age period, populations were relatively small. Now there are large populations. Does population size have any effect in the speed the population is evolving? Translating the idea into genetic algorithms, a third variable for the tests was derived: population size ( $m$ ). A relatively large population size is considered, such as 1,024 individuals and then divide that population by 16, which makes 64 individuals, for a second population size. The number of generations it takes for the population to reach a 70% ( $g = 0.7$ ) having the same solution is considered. This variable is taken into account for both scenarios (genes in series and genes in columns) and for all mutation rates (5%, 10%, 20%, and 50%).

For testing, a medium-size problem is used. Twenty-seven of the most important cities in México are being considered. The number each city has been assigned and the names of the cities are given in Table 2.

Table 2. List of cities.

Number	City
1	Acapulco
2	Aguascalientes
3	Cancún
4	Cuernavaca
5	Chihuahua
6	Durango
7	Guadalajara
8	Hermosillo
9	León
10	Manzanillo
11	Mérida
12	México City
13	Morelia
14	Monterrey
15	Nuevo Laredo
16	Oaxaca
17	Pachuca
18	Puebla
19	Querétaro
20	San Luis Potosí
21	Tampico
22	Tijuana
23	Toluca
24	Torreón
25	Veracruz
26	Villahermosa
27	Zacatecas

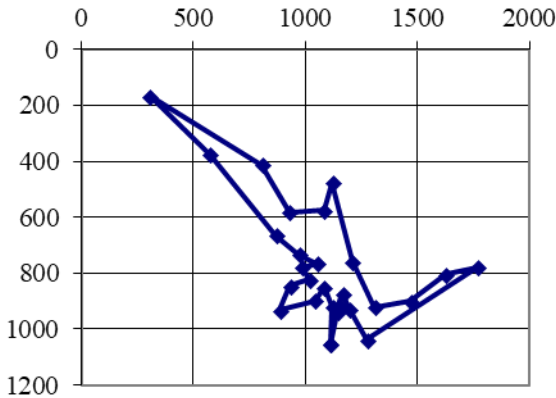
The distances between cities are given in the Appendix.

This problem using linear programming and LINDO/LINGO was solved. The optimal solution to the problem is: 1-23-19-13-10-7-9-2-20-27-6-8-22-5-24-14-15-21-25-26-11-3-16-18-17-12-4. The total distance for this solution is 12,216 km. Notice that this solution is cyclical, and after the last city (4) comes the first city again (1). The fact that the solution is cyclical also means there are  $27 \times 2 = 54$

alternative ways of showing the same solution.

Positions for the cities were obtained in a map. Figure 4 shows the optimal route on a virtual map of México (coordinates are in pixels). In theory, the optimal solution should have no crossings. It can be seen in the map of Figure 4 that there are no crossings.

Figure 4. Route for the optimal solution.



In order to get an idea of the matching between the routes in Figure 4 and an actual map of México and its routes, take a look at Figure 5.

Figure 5. Map of Mexico's routes.



The average of the number of generations it takes to reach 70% agreement in the solution (run or  $r$ ) is calculated. There is also another important matter: how many runs should be considered?

To calculate sample size ( $r$ ), the formula for sample size known as the standard deviation ( $\sigma$ ) was used. The formula was taken from a book by Kvanli, Guynes and Pavur [14] and it is detailed in equation (2).

$$r = \left( \frac{Z_{\alpha/2} \sigma}{E} \right)^2 \quad (2)$$

$Z_{\alpha/2}$  is the value of the normal distribution for a given confidence interval ( $\alpha$ ). In this case,  $Z_{\alpha/2}$  is equal to 1.96, which implies a confidence of 95%. The value for  $\sigma$  is the standard deviation of the population, which typically is approximated using the standard deviation of the sample. Since there is no such value, it is possible to approximate it using equation (3).

$$\sigma \approx \frac{H-L}{4} \quad (3)$$

The reason for the formula in equation (3) is that it is known as an empirical rule that 95.4% of the population will be between  $\mu-2\sigma$  and  $\mu+2\sigma$ . Let  $L = \mu-2\sigma$  be the minimum value expected in the sample and  $H = \mu+2\sigma$  be the maximum value expected, so that  $H-L = 4\sigma$ . Since the statistic indicates the number of runs required, the maximum value for such statistic would be  $r$  ( $H = r$ ) and the minimum value would be zero ( $L = 0$ ). Substituting  $H$  and  $L$  into equation (3) yields equation (4).

$$\sigma \approx \frac{r}{4} \quad (4)$$

Substituting  $\sigma$  from equation (4) into equation (2) yields equation (5).

$$r = \left( \frac{Z_{\alpha/2} r}{4E} \right)^2 \quad (5)$$

Solving for  $r$  from equation (5) yields equation (6).

$$r = \left( \frac{4E}{Z_{\alpha/2}} \right)^2 \quad (6)$$

The only parameter that is not being considered is  $E$ , which is the absolute error. A value for  $E$  of 7 is used. Applying equation (6) yields the sample size: 200, as indicated in equation (7).

$$r = \left( \frac{4 \times 7}{1.96} \right)^2 = 204.08 \approx 200 \quad (7)$$

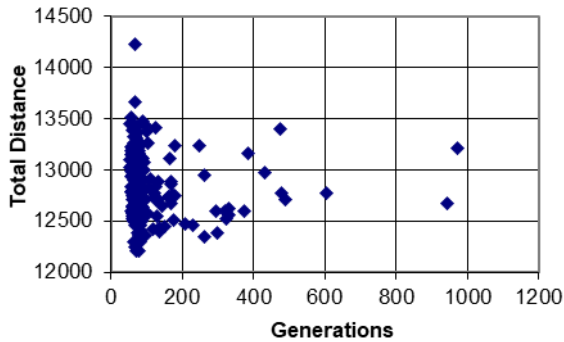
## V. RESULTS

What is a run? There are five steps in genetic algorithms as seen in section 3.B. These steps are: 1. Encoding/Genesis, 2. Evaluation, 3. Crossover, 4. Mutation, and 5. Decoding. After creating a population (step 1), the fitness of all individuals in such population is evaluated (step 2), then comes to reproduce them according to the scheme specific to such sequence (step 3) and some of the individuals mutate according to a given mutation rate (step 4). Then, steps 2, 3, and 4 are iterated a number of times ( $i$ ) until there is at least  $g \times 100\%$  of individuals with the same solution. A run is precisely such sequence of repetitions for a given population. After there is agreement on the solution, the number of generations it took to reach such result, the total distance for the given solution string, as well as the result itself (the sequence of cities to visit) are recorded in files called, by default, output1.txt and output2.txt (step 5). A total of 200 runs are used and the resulting data is recorded. The average of the numbers of generations and total distances for all 200 runs is calculated.

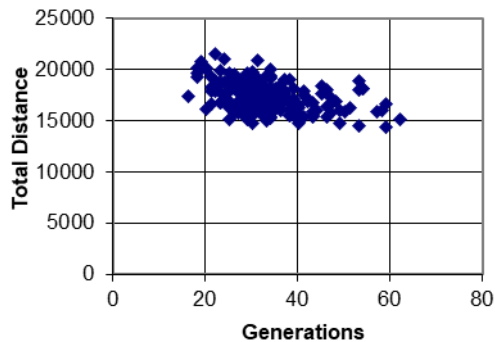
The first set of 200 runs is for a mutation rate of 5% ( $w=0.05$ ) and a population size of 1,024 ( $m=1,024$ ) using genes in series. Figure 6a shows an XY scatterplot of the data. Right next to it, in Figure 6b, is the same mutation rate but for a population of 64 ( $m = 64$ ).

Figure 6. 200 runs for  $w = 0.05$  having genes in series.

a.  $m = 1,024$



b.  $m = 64$

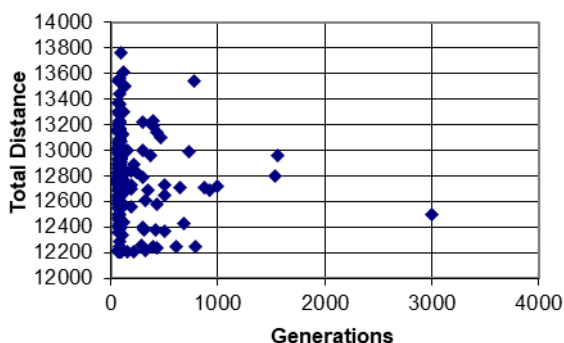


The best run for  $m = 1,024$  generated a route with the same (minimum) total distance of 12,216 km. The route obtained was 19-13-7-10-9-2-20-27-6-8-22-5-24-15-14-21-25-26-11-3-16-18-17-12-4-1-23. This route is the same as the minimum total distance route obtained using linear programming, except for the pairs 7-10 and 15-14, which are reversed. Notice that the linear programming solution starts with 1-23-19-13..., whereas this solution starts with 19-13... Both are the same except for the two pairs mentioned. Nevertheless, both have the same minimum total distance of 12,216 km.

Figure 7 and Figure 8 have the same results for mutation rates of 10% and 20%, respectively. As expected, there could not be obtained results for a mutation rate of 50% given the fact that  $g+w$  would be equal to 1.2, which is greater than 1; see equation (1).

Figure 7. 200 runs for  $w = 0.10$  having genes in series.

a.  $m = 1,024$



b.  $m = 64$

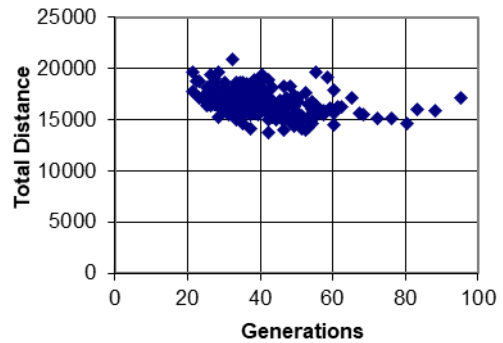
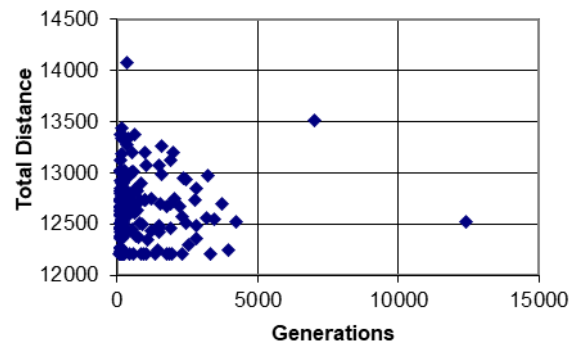
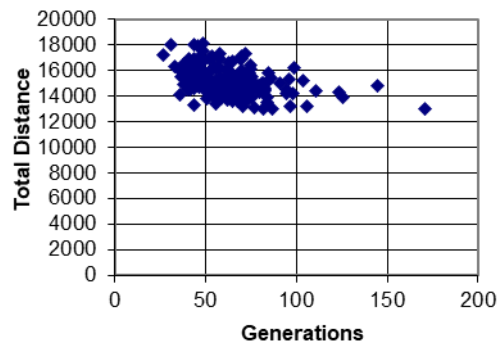


Figure 8. 200 runs for  $w = 0.20$  having genes in series.

a.  $m = 1,024$



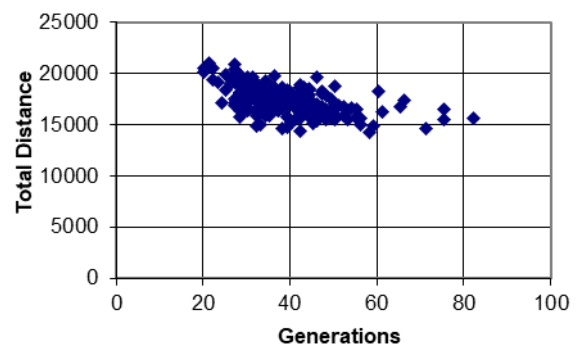
b.  $m = 64$



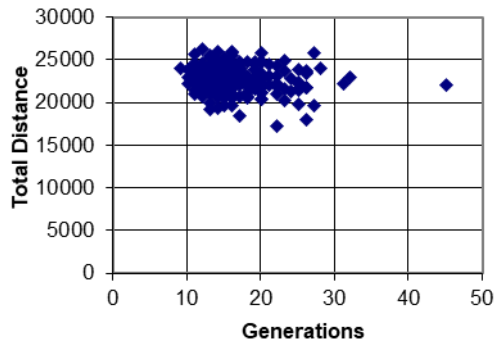
The same set of results was obtained for genes in columns, as shown in Figure 9, Figure 10, and Figure 11 for mutation rates of 5%, 10%, and 20%, respectively. Once again there could not be results for a mutation rate of 50%.

Figure 9. 200 runs for  $w = 0.05$  having genes in columns.

a.  $m = 1,024$



b.  $m = 64$



b.  $m = 64$

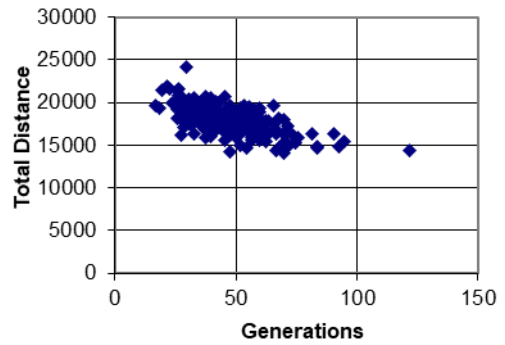
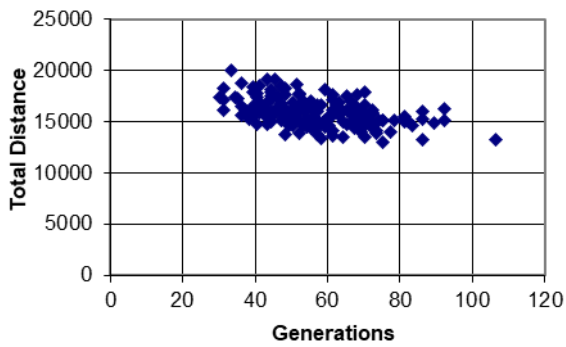


Figure 10. 200 runs for  $w = 0.10$  having genes in columns.  
a.  $m = 1,024$



b.  $m = 64$

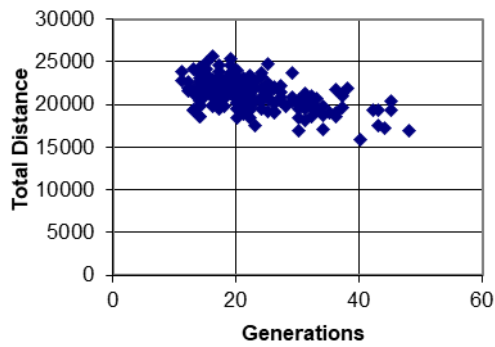
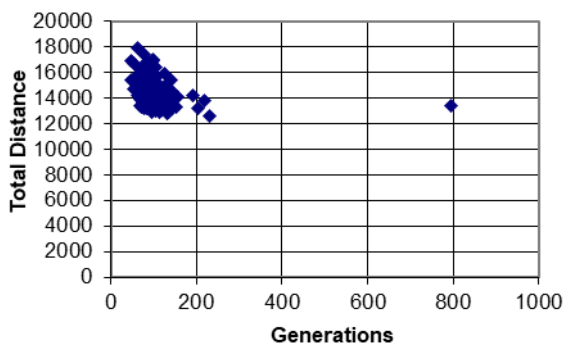
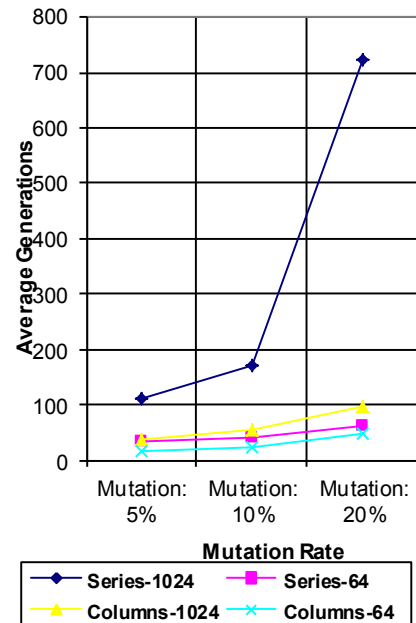


Figure 11. 200 runs for  $w = 0.20$  having genes in columns.  
a.  $m = 1,024$

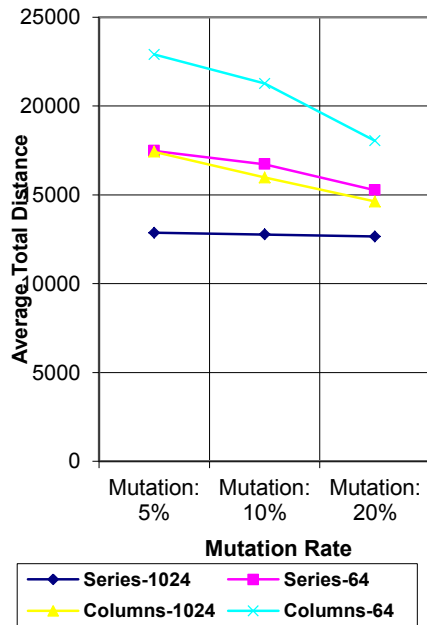


The average number of generations of each set of 200 runs was also calculated, which is shown in Figure 12a, as well as the average total distance of each set of 200 runs, shown in Figure 12b.

Figure 12. Average numbers for each set of 200 runs.  
a. Number of generations



**b. Total distance**



**VI. DISCUSSION AND CONCLUSION**

From Figure 6 it can be seen that it takes less generations (usually about 40) to converge when the population is small ( $m = 64$ ) as opposed to a large population ( $m = 1,024$ ). However, the results obtained are worse when the population is small, which makes sense because a larger population allows the system to consider a wider set of alternatives.

The same can be said for higher mutation rates. Notice that for a mutation of 10% ( $w = 0.10$ ) and for a mutation of 20% ( $w = 0.20$ ) several optimal solutions were found: 6 for 10% mutation and 17 for 20% mutation. Also notice that mutation tends to increase the number of generations it takes to reach a solution. Apparently, a higher mutation rate tends to improve the solution, as long as the mutation is not too high to break the rule from equation (1).

For the case of genes in column, none of the 200 runs of each type of mutation and population size achieved optimal results. Comparing Figure 6 and Figure 9 it can be seen that it takes genes in columns a considerably shorter number of generations to reach agreement on the solution. About 40 generations were required to reach a solution for genes in columns having a mutation rate of 5% and a large population size of 1,024, whereas it took about 100 for genes in series. Apparently, reproducing the population with genes in columns reaches a solution faster. Also, using genes in columns seldom yields an optimal solution.

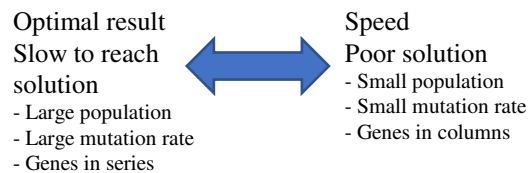
Figure 12a summarizes the results by considering the average of all the figures (total number of generations) for each case: genes in series with  $m = 1,024$ , genes in series with  $m = 64$ , genes in columns with  $m = 1,024$  and genes in columns with  $m = 64$ . It can be seen that the case with the lowest average number of generations to reach a solution is genes in columns with a small population size, followed by genes in series with a small population size. Next is genes in columns with a large population size and finally genes in series with a large population size. Clearly, if the idea is to reach a solution fast there should be used a small population

size. Also, as the mutation rate increases, the average number of generations to reach a solution increases too. The extreme example is the case of genes in series, a large population ( $m = 1,024$ ) and a high mutation rate ( $w = 0.20$ ), exceeding 700 generations on average.

Figure 12b shows the same four cases applied to the average total distance of the resulting tour obtained. The case with the minimum averaged total distance is series with a large population closely followed by columns with a large population. The worst cases were series with small population and finally columns with small population. Thus, in order to achieve a better result (minimum total distance of the tour), it is best to use a large population size. Notice also that as the mutation rate increases, the results are better for all four cases.

In short, a small population size increases speed but reduces performance. Also, a large mutation rate increases performance. Working with genes in series increases performance, whereas genes in columns reduces the number of runs required to reach a solution. For optimal performance (better result), it is recommended to use a large population, large mutation rate and genes in series. The mutation rate should be 10% or 5% under the largest value allowed according to equation (1). In the case under study, the population limit ( $g$ ) is 70%. Thus, the best mutation rate should be  $100\%-70\%-10\% = 20\%$  or maybe even  $100\%-70\%-5\% = 25\%$ . For optimal speed (faster to the result), it is recommended to use a small population, small mutation rate, and genes in columns. Figure 13 summarizes the results obtained.

**Figure 13. Summary of findings.**



There is a tradeoff between the population size, the mutation rate, and the chromosome portraying method. If good performance (optimal solution) is the objective, a large population and large mutation rate, but not as large as to break equation (1), should be considered. This of course results in larger solution times, so a balance among the variables should be attained. If a fast result is the objective, a small population size and a small mutation rate should be used, but not so small as to render useless solutions.

Once the target percentage of the population with the same result ( $g$ ) has been decided, the best mutation rate ( $w$ ) is simply between 1% and 5% below  $1-g$ ; that is between  $w=1-g-0.01$  and  $w=1-g-0.05$ .



APPENDIX. DISTANCES AMONG CITIES

	14	13	12	11	10	9	8	7	6	5	4	3	2	1	Destination/ Origin
1313	1313	652	388	1679	675	781	2337	934	1308	1856	299	1994	899	100000	1
588	588	322	511	1843	549	128	1654	250	433	979	600	2158	100000	899	2
2378	2378	1953	1651	315	2442	2044	3600	2197	2571	3119	1706	100000	2158	1994	3
1014	1014	391	89	1387	880	482	2038	635	1009	1557	100000	1706	600	299	4
783	783	1301	1468	2800	1476	1107	690	1177	667	100000	1557	3119	979	1856	5
563	563	755	920	2252	930	561	1221	631	100000	667	1009	2571	433	1308	6
788	788	268	545	1878	299	230	1403	100000	631	1177	635	2197	250	934	7
1473	1473	1671	1949	3281	1630	1633	100000	1403	1221	690	2038	3600	1654	2337	8
714	714	194	393	1725	530	100000	1633	230	561	1107	482	2044	128	781	9
1085	1085	484	791	2123	100000	530	1630	299	930	1476	880	2442	549	675	10
2063	2063	1634	1332	100000	2123	1725	3281	1878	2252	2800	1387	315	1843	1679	11
925	925	302	100000	1332	791	393	1949	545	920	1468	89	1651	511	388	12
903	903	100000	302	1634	484	194	1671	268	755	1301	391	1953	322	652	13
100000	100000	903	925	2063	1085	714	1473	788	563	783	1014	2378	588	1313	14
224	224	1138	1148	2287	1309	938	1697	1010	787	1007	1238	2602	812	1537	15
1359	1359	756	454	1179	1245	847	2403	1000	1374	1922	456	1494	965	703	16
900	900	390	88	1391	879	410	1995	592	933	1482	177	1710	528	476	17
1048	1048	425	123	1209	914	516	2072	669	1043	1591	178	1528	634	477	18
708	708	195	215	1547	592	178	1763	360	701	1247	304	1866	296	603	19
517	517	397	423	1755	635	197	1720	336	499	1045	512	2074	168	811	20
530	530	790	488	1533	1041	603	2003	742	905	1313	577	1848	574	876	21
2338	2338	2536	2814	4148	2495	2498	865	2268	2086	1555	2903	4465	2519	3202	22
903	903	238	64	1398	727	373	1885	482	896	1442	153	1715	491	285	23
316	316	838	998	2326	1013	644	1157	714	247	467	1085	2647	518	1384	24
1023	1023	720	418	1040	1209	811	2367	964	1338	1898	476	1359	929	775	25
1504	1504	1075	773	559	1564	1166	2722	1319	1893	2241	828	878	1284	1120	26
458	458	452	617	1949	527	258	1524	328	303	849	706	2268	130	1005	27

27	26	25	24	23	22	21	20	19	18	17	16	15	Destination/ Origin
1005	1120	775	1384	285	3202	876	811	603	477	476	703	1537	1
130	1284	929	518	491	2519	574	168	296	634	528	965	812	2
2268	878	1359	2647	1715	4465	1848	2074	1866	1528	1710	1494	2602	3
706	828	476	1085	153	2903	577	512	304	178	177	456	1238	4
849	2241	1898	467	1442	1555	1313	1045	1247	1591	1482	1922	1007	5
303	1893	1338	247	896	2086	905	499	701	1043	933	1374	787	6
328	1319	964	714	482	2268	742	336	360	669	592	1000	1010	7
1524	2722	2367	1157	1885	865	2003	1720	1763	2072	1995	2403	1697	8
258	1166	811	644	373	2498	603	197	178	516	410	847	938	9
527	1564	1209	1013	727	2495	1041	635	592	914	879	1245	1309	10
1949	559	1040	2326	1398	4148	1533	1755	1547	1209	1391	1179	2287	11
617	773	418	998	64	2814	488	423	215	123	88	454	1148	12
452	1075	720	838	238	2536	790	397	195	425	390	756	1138	13
458	1504	1023	316	903	2338	530	517	708	1048	900	1359	224	14
682	1728	1247	540	1127	2582	754	741	932	1272	1124	1583	100000	15
1071	620	382	1450	518	3266	829	877	669	331	513	100000	1583	16
630	832	422	1020	152	2880	387	438	232	182	100000	513	1124	17
740	650	298	1119	187	2937	496	546	338	100000	182	331	1272	18
398	988	633	775	195	2626	499	202	100000	338	232	669	932	19
196	1198	841	582	408	2585	406	100000	202	546	438	877	741	20
502	974	493	846	552	2868	100000	406	499	496	387	829	754	21
2389	3587	3232	2022	2750	100000	2868	2585	2626	2937	2880	3266	2582	22
593	837	482	970	100000	2750	552	408	195	187	152	518	1127	23
386	1769	1414	100000	970	2022	846	582	775	1119	1020	1450	540	24
1035	481	100000	1414	482	3232	493	841	633	298	422	382	1247	25
1390	100000	481	1769	837	3587	974	1198	988	650	832	620	1728	26
100000	1390	1035	386	593	2389	502	196	398	740	630	1071	682	27

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